

Completions, K-categories and Commutative Probability Monads
To DANA SCOTT on his 90th Birthday

Michael Mislove
Tulane University

International Symposium on Domain Theory
Singapore and AWOE, July 4 - 6, 2022

Joint work with Xiaodong Jia, Bert Lindenhovius and Vladimir Zamdzhiev
Supported by US AFOSR

Some History

- 1972 *Continuous lattices* by D. S. Scott, in which the injective T_0 -spaces X are characterized as retracts of $\mathbb{S}^{\mathcal{O}(X)}$ under maps preserving directed suprema, where \mathbb{S} denotes *Sierpinski space*.
- Implies continuous lattices are *sober spaces* in the Scott topology.

Some History

- 1972 *Continuous lattices* by D. S. Scott, in which the injective T_0 -spaces X are characterized as retracts of $\mathbb{S}^{\mathcal{O}(X)}$ under maps preserving directed suprema, where \mathbb{S} denotes *Sierpinski space*.
- 1989 (M-) The *sobrification* of an algebraic poset is an algebraic domain which forms the DCPO-completion of the underlying poset.
- 1992 (H. Zhang) Same result holds for *continuous posets*.

Some History

- 1972 *Continuous lattices* by D. S. Scott, in which the injective T_0 -spaces X are characterized as retracts of $\mathbb{S}^{O(X)}$ under maps preserving directed suprema, where \mathbb{S} denotes *Sierpinski space*.
- 1989 (M-) The *sobrification* of an algebraic poset is an algebraic domain which forms the DCPO-completion of the underlying poset.
- 1992 (H. Zhang) Same result holds for *continuous posets*.
- 2008 (*DCPO-completion of posets*, Zhao & Fan) Described a DCPO-completion of a poset that is finer than the sobrification.
- 2009 (*D-completions and the d-topology*, Keimel and Lawson) Gave topological account of Zhao & Fan's construction, and introduced K -categories, which provide further examples.
- In this talk, we'll apply these results to produce new *commutative* probabilistic monads over DCPO.

We begin by explaining the problem that inspired our work.

What's the Problem?

- The application of interest is *statistical programming*: functional languages that sample from probability distributions.

These are viewed as the simply-typed lambda calculus extended with a *probability monad*.

What's the Problem?

- The application of interest is *statistical programming*: functional languages that sample from probability distributions.

These are viewed as the simply-typed lambda calculus extended with a *probability monad*.

- *Natural model*: Sub-probabilistic power domain $\mathcal{V}_{\leq 1}(P)$ over a dcpo P , the family of *Scott-continuous valuations*: maps $\mu: \sigma(P) \rightarrow [0, 1]$ satisfying:

- $\mu(\emptyset) = 0$
- $\mu(U \cup V) + \mu(U \cap V) = \mu(U) + \mu(V)$
- $\mu(\bigcup_{U \in D} U) = \sup_{U \in D} \mu(U)$, if $D \subseteq \sigma(P)$ is a directed family.

What's the Problem?

- The application of interest is *statistical programming*: functional languages that sample from probability distributions.

These are viewed as the simply-typed lambda calculus extended with a *probability monad*.

- *Natural model*: Sub-probabilistic power domain $\mathcal{V}_{\leq 1}(P)$ over a dcpo P , the family of *Scott-continuous valuations*: maps $\mu: \sigma(P) \rightarrow [0, 1]$
- $\mathcal{V}_{\leq 1}$ defines a monad over DOM – category of domains and Scott continuous maps. In fact, $\mathcal{V}_{\leq 1}$ is *commutative*; i.e., the Fubini-like equation

$$\int_{x \in P} \int_{y \in Q} \chi_U(x, y) d\mu d\nu = \int_{y \in Q} \int_{x \in P} \chi_U(x, y) d\nu d\mu,$$

holds, for domains P, Q , valuations $\mu \in \mathcal{V}_{\leq 1}(P), \nu \in \mathcal{V}_{\leq 1}(Q)$ and $U \subseteq P \times Q$ Scott open. *But*, DOM is not Cartesian closed.

- *Jung-Tix Problem*: There is no known Cartesian closed category of domains on which $\mathcal{V}_{\leq 1}$ defines a monad.

What's the Problem?

- The application of interest is *statistical programming*: functional languages that sample from probability distributions.

These are viewed as the simply-typed lambda calculus extended with a *probability monad*.

- *Natural model*: Sub-probabilistic power domain $\mathcal{V}_{\leq 1}(P)$ over a dcpo P , the family of *Scott-continuous valuations*: maps $\mu: \sigma(P) \rightarrow [0, 1]$
- However, $\mathcal{V}_{\leq 1}$ also defines a monad on DCPO, the category of dcpos and Scott continuous maps, which is Cartesian closed.
- *But*: $\mathcal{V}_{\leq 1}(P)$ is not known to be commutative over DCPO, so there is no proof that Fubini holds.
- So we search for *submonads* of $\mathcal{V}_{\leq 1}$ that are commutative over DCPO.
 - Xiaodong will describe some of the monads we found in his talk on Wednesday.
 - I want to outline the mathematical results that underpin such monads.

The DCPO-completion of a Poset

- *Zhao & Fan*: Defined the *d-topology* on a poset and the associated *d-completion*.¹
 - A subset $A \subseteq P$ is *d-closed* if A is closed under existing suprema of directed sets: i.e., if $D \subseteq A$ is directed and $\sup D \in P$ exists, then $\sup D \in A$.
 - This defines the closed sets of the *d-topology*: the union of finitely many d-closed sets is d-closed (by short argument), and any intersection of d-closed sets is d-closed.
 - Scott-closed subsets are d-closed, so the d-topology refines the Scott topology.
 - The d-closed subsets of a dcpo are exactly the sub-dcpo.
 - Any lower set is d-open, so $\downarrow x$ is d-clopen for each $x \in P$.

¹Zhao & Fan use the notation D-topology, etc., but that clashes with further results we discuss next.

The DCPO-completion of a Poset

- *Zhao & Fan*: Defined the *d-topology* on a poset and the associated *d-completion*.
- A dcpo Q is a *d-completion* of a poset P if there is $\eta: P \rightarrow Q$ Scott continuous satisfying

$$\begin{array}{ccc} Q & \overset{\exists \bar{f}}{\dashrightarrow} & R \\ \eta \uparrow & \nearrow \forall f & \\ P & & \end{array}$$

where R is a dcpo and f, \bar{f} are Scott continuous.

Any two d-completions of P are isomorphic. Denote this by \bar{P} .

- Can be formed as follows:
 - 1) Embed $(P, \sigma(P))$ in a dcpo Q . E.g., take $Q = \Gamma(P)$.
 - 2) Take the intersection of all sub-dcpo's of Q containing P . This is \bar{P} .

Theorem: \bar{P} is the smallest dcpo satisfying $P \hookrightarrow \bar{P}$ is an embedding in the Scott topology.

Monotone Convergence Spaces

2009 *D-completions and the d-topology*, Keimel and Lawson.

The d-topology is not order-theoretic: $Id: \mathbb{S} \rightarrow \{0, 1\}^b$ is d-continuous.

A *monotone convergence space* is a T_0 -space $(X, \mathcal{O}(X))$ in which each directed subset $D \subseteq X$ in the specialization order ($x \leq_s y$ iff $x \in \overline{\{y\}}$) converges to its supremum, $\sup D$.

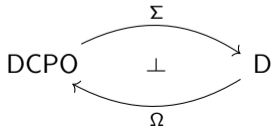
Equivalently, $\Omega(X) = (X, \leq_s)$ is a dcpo and $\mathcal{O}(X) \subseteq \sigma(X, \leq_s)$.

Examples: 1) Any sober space.

2) $\Sigma(P) = (P, \sigma(P))$ – Scott space of dcpo P .

Initially studied as *d-spaces* by Wyler (1981) and later by Ershov (1999).

D - category of monotone convergence spaces and continuous maps.



Monotone Convergence Spaces

2009 *D-completions and the d-topology*, Keimel and Lawson.

D-completion of a T_0 -space X : a d-dense embedding $X \hookrightarrow \tilde{X}$ into a monotone convergence space \tilde{X} . They exist because sober spaces are monotone convergence spaces. (Topological version of Zhao & Fan's d-completion.)

- *D-completions are universal*: for every $\eta: X \hookrightarrow \tilde{X}$, for every $f: X \rightarrow Y (Y \in D)$:

$$\begin{array}{ccc} \tilde{X} & \overset{\exists! \bar{f}}{\dashrightarrow} & Y \\ \eta \uparrow & \nearrow \forall f & \\ X & & \end{array}$$

This implies any two D-completions are homeomorphic. D_c denotes the D-completion.

Monotone Convergence Spaces

2009 *D-completions and the d-topology*, Keimel and Lawson.

- D-completions are universal*: for every $\eta: X \hookrightarrow \tilde{X}$, for every $f: X \rightarrow Y (Y \in D)$:

$$\begin{array}{ccc}
 \tilde{X} & \overset{\exists! \bar{f}}{\dashrightarrow} & Y \\
 \eta \uparrow & \nearrow \forall f & \\
 X & &
 \end{array}$$

This implies any two D -completions are homeomorphic. D_c denotes the D -completion.

- D is a full reflective subcategory of TOP_0 – T_0 -spaces and continuous maps:

$$\begin{array}{ccc}
 & D_c & \\
 \text{TOP}_0 & \begin{array}{c} \curvearrowright \\ \perp \\ \curvearrowleft \end{array} & D \\
 & incl &
 \end{array}$$

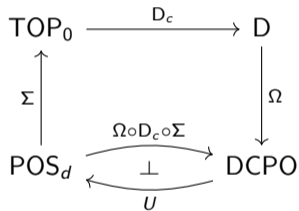
D_c defines an idempotent monad on TOP_0 :

$$f: X \rightarrow Y \quad \mapsto \quad D_c(f) = \overline{\eta_Y \circ f}: D_c(X) \rightarrow D_c(Y)$$

Monotone Convergence Spaces

2009 *D-completions and the d-topology*, Keimel and Lawson.

- The d-completion of a poset in D ; POS_d - posets and Scott continuous maps.



So, $\Omega \circ D_c \circ \Sigma(P)$ is the d-completion of a poset P in the Scott topology.

K-categories

2009 *D-completions and the d-topology*, Keimel and Lawson.

- K-category: a subcategory K of TOP_0 (objects are called “K spaces”) satisfying
 - SOB is a subcategory of K
 - K is closed under homeomorphic images
 - If X is a sober space, the intersection of any family of K -subspaces of X is a K -space
 - If $f: X \rightarrow Y$ in SOB, then f is K -continuous:
i.e., if $Z \subseteq Y$ is a K -subspace, then $f^{-1}(Z)$ is a K -subspace of X ;
equivalently, if $W \subseteq X$ is a subspace, then $f(cl_K(W)) \subseteq cl_K(f(W))$.

Example: D is a K -category.

K-categories

2009 *D-completions and the d-topology*, Keimel and Lawson.

- K-category: a subcategory K of TOP_0 (objects are called “K spaces”) satisfying
 - SOB is a subcategory of K
 - K is closed under homeomorphic images
 - If X is a sober space, the intersection of any family of K -subspaces of X is a K -space
 - If $f: X \rightarrow Y$ in SOB, then f is K -continuous:
i.e., if $Z \subseteq Y$ is a K -subspace, then $f^{-1}(Z)$ is a K -subspace of X ;
equivalently, if $W \subseteq X$ is a subspace, then $f(cl_K(W)) \subseteq cl_K(f(W))$.

Example: D is a K -category.

We can form a K -completion $K_c(X)$ of any T_0 -space: use embedding $\eta: X \hookrightarrow X^s$ into its sobrification and then take

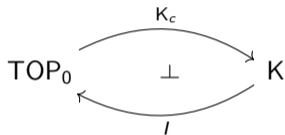
$$\eta_X: X \hookrightarrow K_c(X) = \bigcap \{X' \subseteq X^s \mid X \subseteq X' \text{ a } K \text{ space}\}.$$

The conditions assure that $K_c(X)$ is a K -space and any continuous $f: X \rightarrow Y$ into a K -space admits a unique $\bar{f}: K_c(X) \rightarrow Y$ with $\bar{f} \circ \eta_X = f$

K-categories

2009 *D-completions and the d-topology*, Keimel and Lawson.

Theorem: Each K-category is a full reflective subcategory of TOP_0 :



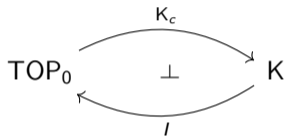
where K_c is the K-completion functor; this defines an idempotent monad on TOP_0 :

$$f: X \rightarrow Y \quad \mapsto \quad K_c(f) = \overline{\eta_Y \circ f}: K_c(X) \rightarrow K_c(Y)$$

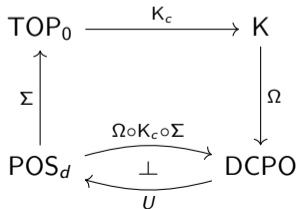
K-categories

2009 *D-completions and the d-topology*, Keimel and Lawson.

Theorem: Each K -category is a full reflective subcategory of TOP_0 :



- If $K \subseteq D$, then $\Omega \circ K_c \circ \Sigma(P)$ defines the K -completion of a poset P in the Scott topology:



Topological View of Valuations

Definition The (extended) probabilistic power domain $\mathcal{V}(X)$ over a topological space $(X, \mathcal{O}(X))$, is the family of maps $\mu: \mathcal{O}(X) \rightarrow \overline{\mathbb{R}}_+$ satisfying:

- $\mu(\emptyset) = 0$
- $\mu(U \cup V) + \mu(U \cap V) = \mu(U) + \mu(V)$
- $\mu(\bigcup_{U \in D} U) = \sup_{U \in D} \mu(U)$, if $D \subseteq \mathcal{O}(X)$ is directed.

We endow $\mathcal{V}(X)$ with the *stochastic order*: $\mu \leq \nu$ iff $\mu(U) \leq \nu(U)$ ($\forall U \in \mathcal{O}(X)$).

- $\mathcal{V}(X)$ is a DCPO – in fact, $\mathcal{V}: \text{TOP}_0 \rightarrow \text{DCPO}$ is a functor, where

$$f \in \text{TOP}_0(X, Y) \quad \mapsto \quad \mu \mapsto \mathcal{V}f(\mu) = \mu \circ f^{-1} \in \text{DCPO}(\mathcal{V}(X), \mathcal{V}(Y)).$$

(Jones) $\mathcal{V} \circ \Sigma: \text{DCPO} \rightarrow \text{DCPO}$ defines a strong monad;

$(\mathcal{V} \circ \Sigma)|_{\text{DOM}}: \text{DOM} \rightarrow \text{DOM}$ is a commutative monad.

Topological View of Valuations

Definition The (*extended*) *probabilistic power domain* $\mathcal{V}(X)$ over a topological space $(X, \mathcal{O}(X))$, is the family of maps $\mu: \mathcal{O}(X) \rightarrow \overline{\mathbb{R}}_+$ satisfying:

- $\mu(\emptyset) = 0$
- $\mu(U \cup V) + \mu(U \cap V) = \mu(U) + \mu(V)$
- $\mu(\bigcup_{U \in D} U) = \sup_{U \in D} \mu(U)$, if $D \subseteq \mathcal{O}(X)$ is directed.

We endow $\mathcal{V}(X)$ with the *stochastic order*: $\mu \leq \nu$ iff $\mu(U) \leq \nu(U)$ ($\forall U \in \mathcal{O}(X)$).

$\mathcal{V}_w X$ denotes $\mathcal{V}X$ in the *weak topology*: $[U > r] = \{\mu \mid \mu(U) > r\} = \{\mu \mid \text{eval}(\mu, U) > r\}$ forms a sub-basis, where $\text{eval}: \mathcal{V}X \times \mathcal{O}(X) \rightarrow \mathbb{R}_+$. Heckmann showed $\mathcal{V}_w(X)$ is sober.

Topological View of Valuations

Definition The (extended) probabilistic power domain $\mathcal{V}(X)$ over a topological space $(X, \mathcal{O}(X))$, is the family of maps $\mu: \mathcal{O}(X) \rightarrow \overline{\mathbb{R}}_+$ satisfying:

- $\mu(\emptyset) = 0$
- $\mu(U \cup V) + \mu(U \cap V) = \mu(U) + \mu(V)$
- $\mu(\bigcup_{U \in D} U) = \sup_{U \in D} \mu(U)$, if $D \subseteq \mathcal{O}(X)$ is directed.

(Goubault-Larrecq & Jia) \mathcal{V}_w defines a monad on TOP_0 :

$\mathcal{V}_w f: \mathcal{V}_w X \rightarrow \mathcal{V}_w Y$ is $\mathcal{V}_w f(\mu) = \mu \circ f^{-1}$ is continuous, for $f: X \rightarrow Y$.

The *unit* is the point valuation $x \mapsto \delta_x = \lambda U. \begin{cases} 1 & x \in U, \\ 0 & \text{otherwise.} \end{cases}$

We can define the Choquet-like integral $\int_X h(x) d\mu = \int_0^\infty \mu(h^{-1}(r, \infty]) dr$ (the Riemann integral) for $h: X \rightarrow \overline{\mathbb{R}}_+$ lower semicontinuous and any space X .

Then the *multiplication* $m: \mathcal{V}_w^2 X \rightarrow \mathcal{V}_w X$ is defined as

$$m(\bar{w}) = \lambda U. \int_{\mathcal{V}_w X} \mu(U) d\bar{w} = \lambda U. \int_0^\infty \bar{w}(\text{eval}(-, U)^{-1}(r, \infty]) dr = \lambda U. \int_0^\infty \bar{w}([U > r]) dr.$$

Topological View of Valuations

Definition The (extended) probabilistic power domain $\mathcal{V}(X)$ over a topological space $(X, \mathcal{O}(X))$, is the family of maps $\mu: \mathcal{O}(X) \rightarrow \overline{\mathbb{R}}_+$ satisfying:

- $\mu(\emptyset) = 0$
- $\mu(U \cup V) + \mu(U \cap V) = \mu(U) + \mu(V)$
- $\mu(\bigcup_{U \in D} U) = \sup_{U \in D} \mu(U)$, if $D \subseteq \mathcal{O}(X)$ is directed.

Theorem For any space X , $\mathcal{V}_w X$ has a canonical cone structure making it a locally linear topological cone. Moreover, $\mathcal{V}_w f: \mathcal{V}_w X \rightarrow \mathcal{V}_w Y$ is continuous and linear, for each $f: X \rightarrow Y$.

$+$: $\mathcal{V}_w X \times \mathcal{V}_w X \rightarrow \mathcal{V}_w X$ is $(\mu + \nu)(U) = \mu(U) + \nu(U)$ for each open set $U \subseteq X$;

$r \cdot -$: $\mathcal{V}_w X \rightarrow \mathcal{V}_w X$ is $(r \cdot \mu)(U) = r \cdot \mu(U)$ for each $r \in \mathbb{R}_+$; and

$U \mapsto 0 \in \overline{\mathbb{R}}_+$ is the zero valuation.

The sub-basis $[U > r]$ consists of (open) half-spaces (i.e., convex subsets whose complement also is convex), which is what locally linear means. The last claim is straightforward from the definitions of $\mathcal{V}_w f(\mu) = \mu \circ f^{-1}$ and of the sub-bases $[U > r]$ for $\mathcal{V}_w X$ and $\mathcal{V}_w Y$.

K-completions of Simple Valuations

Simple valuations over X : $\mathcal{V}_s X = \{\sum_{x \in F} r_x \delta_x \mid r_x \in \mathbb{R}_+ \text{ \& } F \subseteq X \text{ finite}\} \subseteq \mathcal{V}_w X$.

\mathcal{V}_s is a submonad of \mathcal{V}_w

- If $f: X \rightarrow Y$, then $\mathcal{V}_s f(\sum_{x \in F} r_x \cdot \delta_x) = \sum_{x \in F} r_x \cdot \delta_{f(x)}$.
- The unit $x \mapsto \delta_x: X \rightarrow \mathcal{V}_w X$ is simple.
- The multiplication $m: \mathcal{V}_w^2 X \rightarrow \mathcal{V}_w X$ restricts to $\mathcal{V}_s X^2 \rightarrow \mathcal{V}_s X$ via

$$m\left(\sum_{x \in F} r_x \left(\sum_{y \in G_x} s_{x,y} \delta_{x,y}\right)\right) = \sum_{x,y} r_x s_{x,y} \delta_{x,y}.$$

K-completions of Simple Valuations

Simple valuations over X : $\mathcal{V}_s X = \{\sum_{x \in F} r_x \delta_x \mid r_x \in \mathbb{R}_+ \text{ \& } F \subseteq X \text{ finite}\} \subseteq \mathcal{V}_w X$.

\mathcal{V}_s is a submonad of \mathcal{V}_w

In particular, there is an embedding $\eta_{\mathcal{V}_s X}: \mathcal{V}_s X \hookrightarrow \mathcal{V}_w X$, which is sober.

If K is a K -category, it follows that the K -completion $\mathcal{V}_K X \stackrel{\text{def}}{=} K_c(\mathcal{V}_s X)$ is a sub-dcpo of $\mathcal{V}_w X$.

Theorem For each K -category K , \mathcal{V}_K is a monad on TOP_0 .

Proof: $\mathcal{V}_K X$ is a subcone of $\mathcal{V}_w X$, so it is a locally linear topological cone, and each continuous linear map $f: \mathcal{V}_s X \rightarrow \mathcal{V}_s Y$ satisfies $\eta_{\mathcal{V}_K} \circ f(\mathcal{V}_s X) \subseteq \mathcal{V}_K X$, so $K_c(f): \mathcal{V}_K X \rightarrow \mathcal{V}_K Y$. The linearity of $K_c(f)$ follows from the density of $\mathcal{V}_s X$ in $\mathcal{V}_K X$ and the continuity of addition on $\mathcal{V}_K Y$. This shows \mathcal{V}_K is an endofunctor on TOP_0 .

The unit is the point valuation $x \mapsto \delta_x$, and if $f: X \rightarrow \mathcal{V}_K Y$ is continuous, then

$$K_c(f) = \mu \mapsto (U \mapsto \int_X f(x)(U) d\mu): \mathcal{V}_K X \rightarrow \mathcal{V}_K Y$$

is well-defined and continuous. This shows \mathcal{V}_K defines a submonad of \mathcal{V}_w .

K-completions of Simple Valuations

Simple valuations over X : $\mathcal{V}_s X = \{\sum_{x \in F} r_x \delta_x \mid r_x \in \mathbb{R}_+ \text{ \& } F \subseteq X \text{ finite}\} \subseteq \mathcal{V}_w X$.

\mathcal{V}_s is a submonad of \mathcal{V}_w

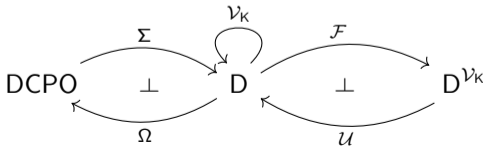
In particular, there is an embedding $\eta_{\mathcal{V}_s X}: \mathcal{V}_s X \hookrightarrow \mathcal{V}_w X$, which is sober.

If K is a K -category, it follows that the K -completion $\mathcal{V}_K X \stackrel{\text{def}}{=} K_c(\mathcal{V}_s X)$ is a sub-dcpo of $\mathcal{V}_w X$.

Theorem For each K -category K , \mathcal{V}_K is a monad on TOP_0 .

Theorem For each K -category $K \subseteq D$, $\mathcal{V}_{K, \leq} \stackrel{\text{def}}{=} \Omega \circ \mathcal{V}_K \circ \Sigma$ is a monad on DCPO .

Proof: \mathcal{V}_K is a monad on TOP_0 , and since K is a full subcategory of D , it follows that \mathcal{V}_K restricts to a monad on D . Writing $\mathcal{V}_K = \mathcal{U} \circ \mathcal{F}$, we have $\Omega \circ \mathcal{V}_K \circ \Sigma = (\Omega \circ \mathcal{U}) \circ (\mathcal{F} \circ \Sigma)$, and adjoints compose:



It's straightforward to see that the unit and multiplication of \mathcal{V}_K are transported to DCPO .

K-completions of Simple Valuations

Simple valuations over X : $\mathcal{V}_s X = \{\sum_{x \in F} r_x \delta_x \mid r_x \in \mathbb{R}_+ \text{ \& } F \subseteq X \text{ finite}\} \subseteq \mathcal{V}_w X$.

\mathcal{V}_s is a submonad of \mathcal{V}_w

Theorem For each K -category $K \subseteq D$, $\mathcal{V}_{K, \leq} \stackrel{\text{def}}{=} \Omega \circ \mathcal{V}_K \circ \Sigma$ is a monad on DCPO.

In particular, this applies to the D -completion $\mathcal{V}_D X$.

Theorem The monad $\mathcal{M} \stackrel{\text{def}}{=} \mathcal{V}_{D, \leq 1}$ is commutative on DCPO.

Proof: If $\mu = \sum_{x \in F} r_x \cdot \delta_x \in \mathcal{M}X$ and $\nu \in \mathcal{M}Y$, then

$$\int_X \int_Y \chi_U(x, y) d\mu d\nu = \sum_{x \in F} r_x \cdot \int_Y \chi_U(x, y) d\nu = \int_Y \sum_{x \in F} r_x \cdot \chi_U(x, y) d\nu = \int_Y \int_X \chi_U(x, y) d\mu d\nu,$$

and since $\mathcal{V}_s X$ is dense in $\mathcal{M}X$ and integration is continuous, the equation holds for all $\mu \in \mathcal{M}X$.

K-completions of Simple Valuations

Simple valuations over X : $\mathcal{V}_s X = \{\sum_{x \in F} r_x \delta_x \mid r_x \in \mathbb{R}_+ \text{ \& } F \subseteq X \text{ finite}\} \subseteq \mathcal{V}_w X$.

\mathcal{V}_s is a submonad of \mathcal{V}_w In particular, this applies to the D-completion $\mathcal{V}_D X$.

Theorem The monad $\mathcal{M} \stackrel{\text{def}}{=} \mathcal{V}_{D, \leq 1}$ is commutative on DCPO.

This also applies to any full K-subcategory K of D; in fact, for such a K, we have

$\mathcal{V}_{s, \leq 1} \subseteq \mathcal{M} \subseteq \mathcal{V}_{K, \leq 1} \subseteq \mathcal{P}$ (but $\mathcal{V}_{s, \leq 1}$ isn't a subcategory of DCPO), where $\mathcal{P} = \mathcal{V}_{\text{SOB}, \leq 1}$.

Heckmann showed that \mathcal{P} consists of *point continuous valuations*; hence the name.

Another example we know of is the full K-subcategory W of D consisting of *well-filtered spaces*. So the commutative monads we know so far are:

$$\mathcal{V}_{s, \leq 1} \subseteq \mathcal{M} \subseteq \mathcal{W} \subseteq \mathcal{P} \subseteq \mathcal{Z}$$

The last – \mathcal{Z} – denotes the category of *central valuations*: those for which integration satisfies the Fubini equation with any valuation in \mathcal{V} for the other component. This category exists by abstract reasoning, and doesn't rely on \mathcal{V}_s being a dense subcategory relative to some completion operation. It also is the only one we know that contains the pushforward of Lebesgue measure by a lower semicontinuous map to a DCPO.

Some References

- Completing Simple Valuations in K -categories, X. Jia and M. Mislove, *Topology and Its Applications*, *in press*, Available online at <https://arxiv.org/abs/2002.01865v2>
- Commutative Monads for Probabilistic Programming Languages, X. Jia, B. Lindenhovius, M. Mislove and V. Zamdzhiev, in: *LICS '21: Proceedings of the 36th Annual ACM/IEEE Symposium on Logic in Computer Science (2021)*
DOI: <https://doi.org/10.1109/LICS52264.2021.9470611>
Also available at <https://arxiv.org/pdf/2102.00510.pdf>
- The Central Valuations Monad (Early Ideas), X. Jia, M. Mislove and V. Zamdzhiev, 9th Conference on Algebra and Coalgebra in Computer Science (CALCO 2021), *Leibniz International Proceedings in Informatics (LIPIcs)*,
DOI: [10.4230/LIPIcs.CALCO.2021.18](https://doi.org/10.4230/LIPIcs.CALCO.2021.18)
Also available at: <https://arxiv.org/abs/2111.10873>